Sample questions for midterm 2
CS 421, Spring 2009

1. Recall that in lectures 11 and 12 we gave several translation schemes for expressions and statements: [e] produces a pair containing the code for e and the location of the result; [S] gives the code for S; $[e]_x$ gives the code to evaluate e and put its value in x; $[e]_{Lt,Lf}$ translates a Boolean expression e to code that branches to either Lt or Lf (this is called the "short-circuit evaluation scheme"), and $[S]_L$ translates S in a context in which L is the label for a break statement.

a. Generate code for the repeat-until statement: "repeat S until e" executes S and tests e, and repeats until e becomes true. Thus, it is equivalent to "S; while !e do S". Do this in two ways: (i) Using the regular scheme [e] to evaluate the condition; and (ii) Using the short-circuit evaluation scheme for e.

b. Generate code for a multiple assignment statement: (x1, x2) = (e1, e2), which does both assignments "in parallel." Note that this is not that same as doing one assignment followed by the other, because variables x1 and x2 may appear in expressions e1 and e2. Use evaluation scheme $[e]_x$ where appropriate.

c. Give two schemes for conditional expression e1 ? e2 : e3, which gives the value of e2 if e1 is true, or e3 if e1 is false. The two schemes you should provide are (a) the standard [e1 ? e2 : e3], and (b) the assignment scheme $[e1 ? e2 :e3]_x$. You should use the short-circuit evaluation scheme for e1.

2. (a) Name the two parts of a compiler's front end.

(b) Name the ~~two~~ three parts of a compiler's back end. 1 gen IR 2 optimize 3 code gen

(c) What are the two outputs of the front end?

3. Name the items in an ~~activation record.~~ stack frame.

4. Give two advantages of the copying garbage collection algorithm over the non-copying (mark-and-sweep) algorithm. 1. Usually faster because only touches reachable cells 2. Keep active data in highest level of hierarchy

5. Give two advantages of the non-copying (mark-and-sweep) garbage collection algorithm over the copying algorithm. 1. Uses entire memory 2. Don't need to change pointers

6. Reference counting is not a popular algorithm. What drawback of this algorithm is the reason? 3. Don't need to move large heap objects. Cycles in heap.

7. In ~~APL, define multmat n which~~ gives an nxn matrix where position i,j has the value i*j.

multmat 4;;
1 2 3 4

Translation schemes

$[e]$ - code for e + loc. of result

$[S]$ - code for S

$[e]_x$ - code to store e in x

$[S]_L$ - code for S within a loop or switch stmt, where L is target of a break.

$[e]_{Lt,Lf}$ - "short-circuit evaluation" for boolean expr's; code to jump to Lt or Lf, depending on value of e.

[repeat S until e]

```
= let (il, loc) = [e]
  let L1, L2 = newlabels()
  in
     L1: [S]
         il
         CJUMP loc, L2, L1
     L2:
```

[repeat S until e]

- Short-circuit version

= ~~let (il, loc) = [e]~~

let L1, L2 = newlabels()

in

L1: [S]

[e]$_{L2, L1}$

L2:

while : while (e)
S

do-while : do S
while (e)

repeat-until : repeat S ≡ do S
until e while (!e)

$[e_1 \; ? \; e_2 : e_3] =$

let $x =$ newlocation()

$L_1, L_2^{L_3} =$ newlabels

$(il_2, x^2) = [e_2]$

$(il_3, x_3) = [e_3]$

in

$[e_1]_{L_1, L_2}$

L_1: il_2

$x = x_2$

JUMP L_3

L_2: il_3

$x = x_3$

L_3:

X

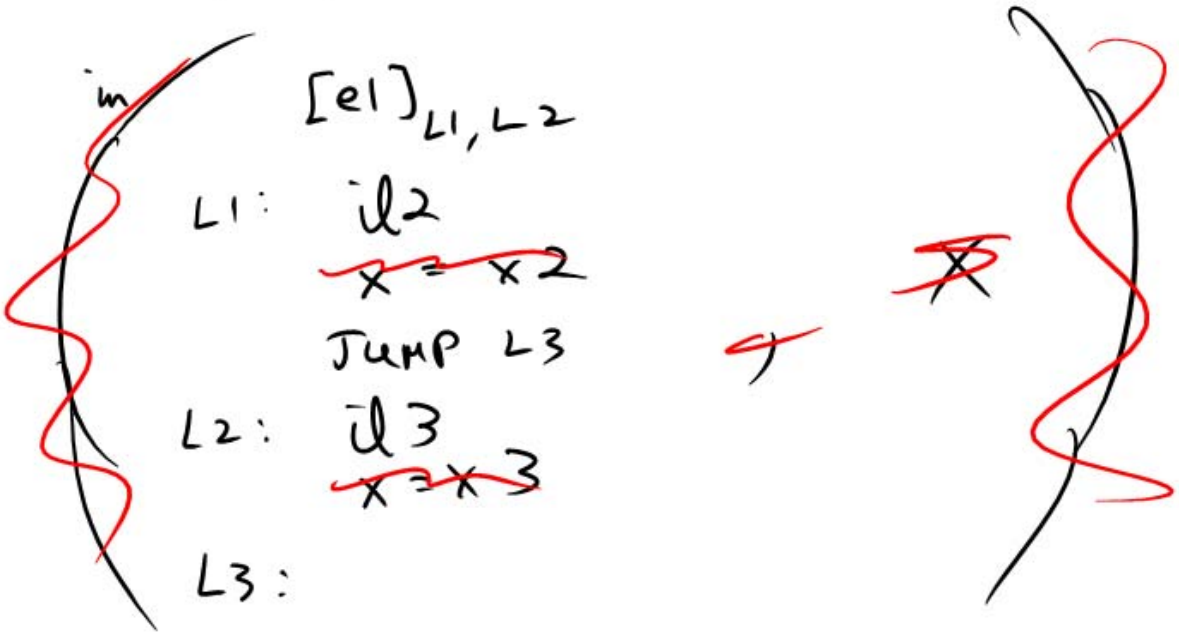
$[e_1 \; ? \; e_2 : e_3]_x =$

let ~~x = newlocation()~~

$L_1, L_2 = \text{newlabels}^{L_3}$

~~(~~ $(il_2, x_2) = [e_2]_x$

~~(~~ $(il_3, x_3) = [e_3]_x$



1 2 3 4
2 4 6 8
3 6 9 12
4 8 12 16

let M = n n p 2n
in (⊙M) × M

1 2 3 4
1 2 3 4
1 2 3 4
1 2 3 4

8. Define the following OCaml functions. *[Exam: we will provide definitions of fold_right and fold_left.]:*

(a) repeat_until: ('a -> bool) -> ('a -> 'a) -> 'a -> 'a. where repeat_until p f x = x, if p x, or f x if p (f x), or f (f x) if p (f (f x)), etc.

(b) sift: ('a -> bool) -> 'a list -> 'a list * 'a list. sift p lis splits lis into a pair of lists (lis1, lis2), with lis1 containing those elements of lis that satisfy p and lis2 the others.

(c) Write sift using fold_right. Specifically, define sift_base and sift_rec so that
fold_right (sift_rec p) lis sift_base = sift p lis

(d) Write an OCaml function that reverses a list, using *fold_right* instead of explicit recursion.

(e) Write a function f such that map f lis returns a list that contains the absolute values of the elements in lis, in the same order. Do not use any library functions in the definition of f.

(f) Using fold_right and no explicit recursion, define a function that concatenates the elements of a string list.

(g) compose_all [f1;f2;…] a = f1 (f2 (… (fn a)…)). Define compose_all *and* say what its type is.

(h) graph_fun f [x1; x2; ...; xn] = [(x1, f x1); (x2, f x2); ...]. Define graph_fun and say what its type is.
`graph_fun`: (α→ β) → α list → (α * β) list, where

9. What does this OCaml program evaluate to:

```
let x = 4
let y = 6
let f = x + y
let x = 8
in f(y+x)
```

x = 4
y = 6
f = fun z → 4 + z
x = 8
f(y+x) = f(6+8) = f 14 = 18

10. Suppose the following Java interface is defined:

```
interface FunObj {
  int apply (int) ;
}
```

A Java class might define a function like this:

```
void map (FunObj f, int[] a) {
```

Compose_all [f1; f2; ...; fn] a =

f1 (f2 (... fn- (fn a) ...))

compose_all : (('α → α) list) → α

→ α

let rec compose_all fl a
= match fl with
[] → a
| f :: fl' → f (compose_all fl' a)

compose_all $[f^1; f^2; \ldots; f^n]$ a =

$f^1\ (f^2\ (\ldots^{for}\ (f^n\ a)\ \ldots))$

fold_right : $(\alpha * \beta \to \beta) \to \alpha$ list $\to \beta \to \beta$

fold_right (fun a b → a b)

this

a

apply

compose_all [succ; decr] 0 = 0

{ fold_right (fun a b → a b) [--] 0

fold_right f $[x_1; \ldots; x_n]$ a = f x_1 (f x_2 (-- f x_n a) --)

→ apply succ (apply pred 0)

(fun a b → a b) succ -1 = (fun a b → a b) pred 0

= succ -1 = 0 = pred 0 = -1

```
graph_fun f [x1; ...; xn] =
  fold_right (fun x a -> (x, f x) :: a)
             l
             []
```

→ $[(x_1, f\,x_1); (x_2, f\,x_2); \ldots]$

```
void map(FunObj f, int[] a) {
    for (int i=0; i<a.length; i++)
        a[i] = f.apply(a[i]);
}
```

a. Define classes **decrobj** and **sqrobj** that implement **FunObj** so that **map (new decrobj(), a, n)** decrements each element of **a**, and **map (new sqrobj(), a, n)** squares each element of **a**.

b. Define a class **compose** that implements **FunObj**:

```
class compose implements FunObj {
    FunObj f, g;
    compose (FunObj f, FunObj g ) {
        this.f = f; this.g = g;
    }

    int apply (int i) {
        return f.apply(g.apply(i));


    }
}
```

that composes function objects, so that, for example, **map(new compose(new sqrobj(), new decrobj()), a, n)** changes every element a[i] to $(a[i]-1)^2$.

```
class decrobj implements FunObj {
    int apply (int n) {
        return n-1;
    }
}

map( new decrobj(), a )
```

Ocaml:

```
compose = fun f g -> fun i -> f (g i)
compose decr sqr
```